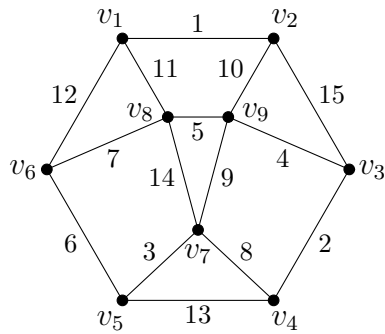


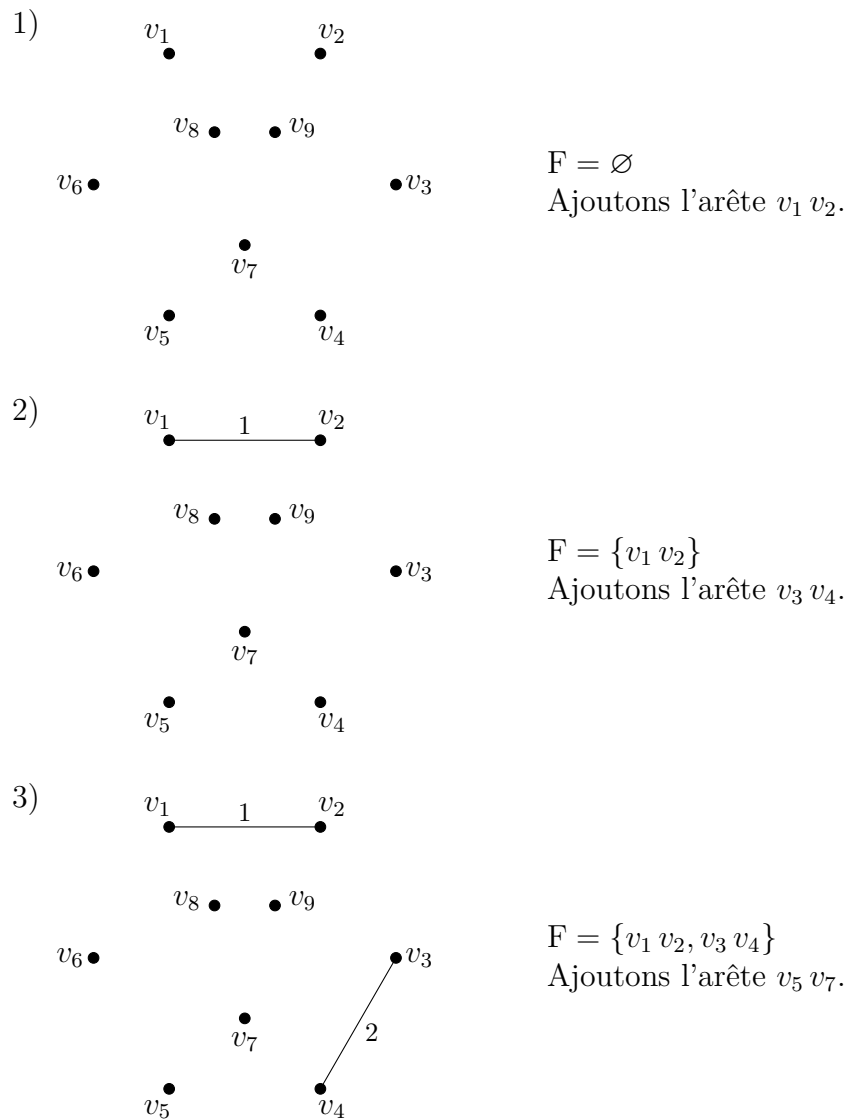
### 3.11 Algorithme de Kruskal

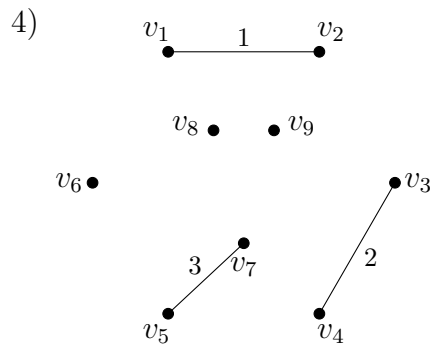


Ordonnons les arêtes du graphe par ordre croissant de poids :

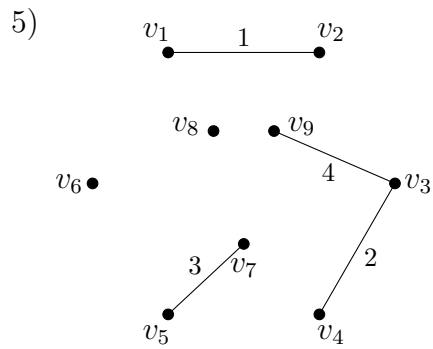
$\underbrace{v_1 v_2}_1, \underbrace{v_3 v_4}_2, \underbrace{v_5 v_7}_3, \underbrace{v_3 v_9}_4, \underbrace{v_8 v_9}_5, \underbrace{v_5 v_6}_6, \underbrace{v_6 v_8}_7, \underbrace{v_4 v_7}_8, \underbrace{v_7 v_9}_9, \underbrace{v_2 v_9}_{10}, \underbrace{v_1 v_8}_{11}, \underbrace{v_1 v_6}_{12}, \underbrace{v_4 v_5}_{13}, \underbrace{v_7 v_8}_{14}, \underbrace{v_2 v_3}_{15}$

Appliquons les différentes étapes de l'algorithme de Kruskal :

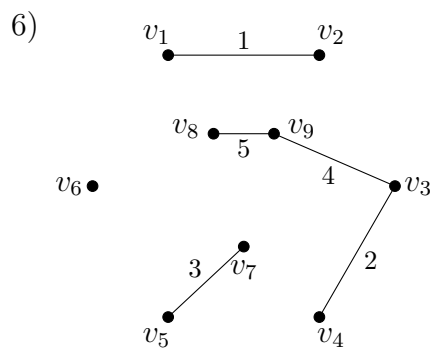




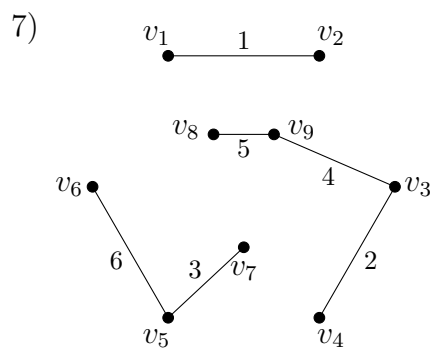
$F = \{v_1 v_2, v_3 v_4, v_5 v_7\}$   
Ajoutons l'arête  $v_3 v_9$ .



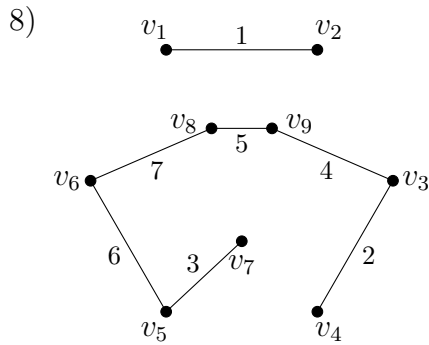
$F = \{v_1 v_2, v_3 v_4, v_5 v_7, v_3 v_9\}$   
Ajoutons l'arête  $v_8 v_9$ .



$F = \{v_1 v_2, v_3 v_4, v_5 v_7, v_3 v_9, v_8 v_9\}$   
Ajoutons l'arête  $v_5 v_6$ .



$F = \{v_1 v_2, v_3 v_4, v_5 v_7, v_3 v_9, v_8 v_9, v_5 v_6\}$   
Ajoutons l'arête  $v_6 v_8$ .

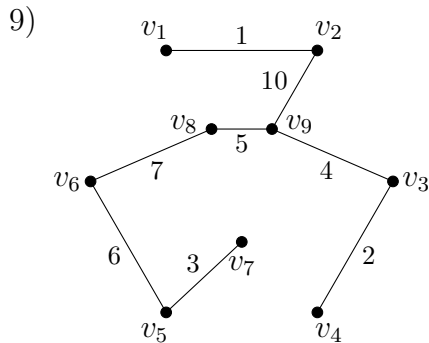


$$F = \{v_1 v_2, v_3 v_4, v_5 v_7, v_3 v_9, v_8 v_9, v_5 v_6, v_6 v_8\}$$

L'ajout de l'arête  $v_4 v_7$  créerait le cycle  $v_3 v_4 v_7 v_5 v_6 v_8 v_9 v_3$ .

L'ajout de l'arête  $v_7 v_9$  produirait le cycle  $v_5 v_6 v_8 v_9 v_7 v_5$ .

Ajoutons l'arête  $v_2 v_9$ .

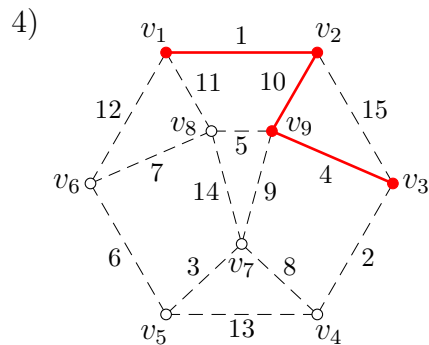
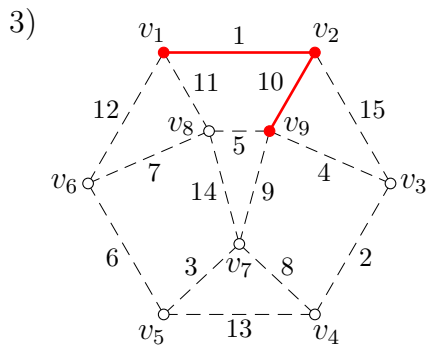
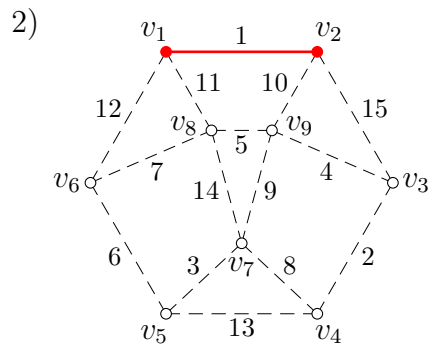
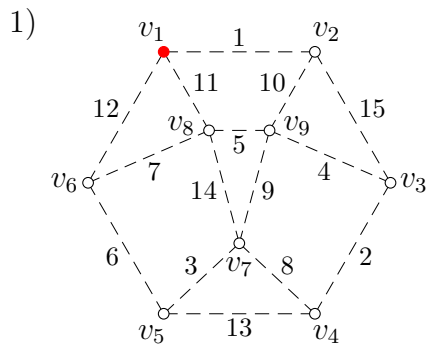


$$F = \{v_1 v_2, v_3 v_4, v_5 v_7, v_3 v_9, v_8 v_9, v_5 v_6, v_6 v_8, v_2 v_9\}$$

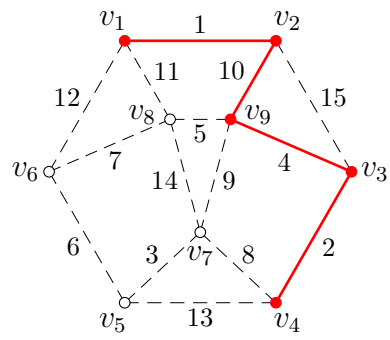
À présent que tous les sommets sont connectés, nous avons obtenu l'arbre de recouvrement de poids minimum.

### Algorithme de Prim

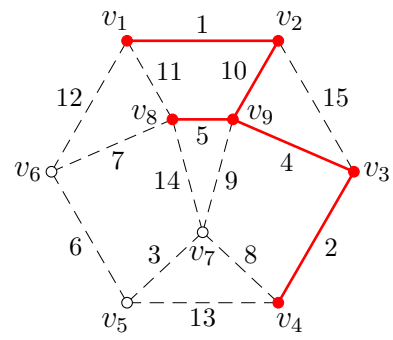
Voici les différentes étapes de l'algorithme de Prim, en marquant initialement le sommet  $v_1$  :



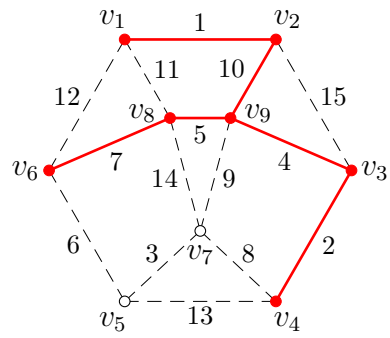
5)



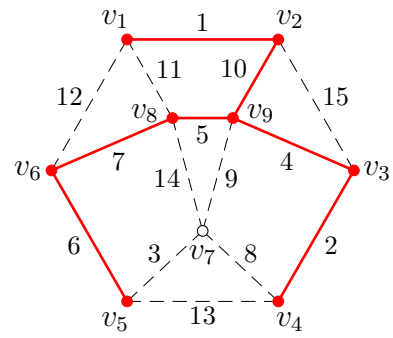
6)



7)



8)



9)

